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## Title of the Invention

Selectable Complexity Turbo Coding System

## Field of the Invention

5        The present invention relates to switchable encoding and decoding of a data sequence, and more particularly but not exclusively to switchable turbo code encoding and decoding.

## Background of the Invention

10        Turbo codes are employed in modern digital communication systems to protect high bit-rate transmitted information from error. Turbo codes are generally constructed as a concatenation of two recursive systematic convolutional codes, linked together by some non-uniform interleaving. The term turbo code originally described the parallel concatenation of two recursive systematic convolutional codes (PCCC). Other possibilities are serial  
15        concatenation (SCCC) and using block codes as component codes. PCCC is now becoming a standard error correction scheme in several wireless air interfaces. For example, the 3GPP (Third Generation Partnership Project for wireless systems) standard employs a PCCC with  $M=8$  states and a block length of up to  $N=5120$  information bits. In a cdma2000 system,  $N$  can be as  
20        high as 20000.

      Although convolutional in nature, turbo codes cannot be decoded directly using a Viterbi decoder. The Viterbi decoding algorithm models a code as a trellis, with branches depicting legal transitions from state to state.

Each state represents a combination of input digits used by the encoder to select the transmitted symbol, and each branch is associated with a branch metric. As each symbol in the received sequence is processed, the Euclidean distance between the received symbol and each possible path through the trellis is measured. A single surviving path is selected for each state.

A trellis diagram corresponding to a turbo code typically has a huge number of states, making implementation of the Viterbi algorithm impractical. Therefore, an iterative approach is employed with two elementary decoders, each associated with one of the two constituent codes. The two decoders are usually serially concatenated, where the first decoder yields weighted, or soft-output decisions that are fed to the second decoder as *a priori* information. The soft-outputs of the second decoder are then fed back to the first decoder for the second iteration, and so on. Only the so-called extrinsic information, i.e. new information that is generated by the decoder, is passed between the decoders.

The optimal soft-output decoder is the so-called MAP (maximum *a posteriori*) decoder, which uses both backward and forward decoding to efficiently determine the soft output. The MAP decoder is optimal in the sense that it minimizes the decoded bit error probability for each information bit based on all received bits. However, because of memory, processing, and numerical tradeoffs, MAP decoding is usually limited to a sub-optimal approximation. Typically, convolutional codes composing a turbo code are graphically represented as a trellis. MAP-type decoders (log-MAP, MAP, max-log-MAP, constant-log-MAP, etc.) utilize forward and backward

generalized Viterbi recursions on the trellis in order to provide soft outputs, as is known in the art.

Because of the Markov nature of the encoded sequence (wherein previous states cannot affect future states or future output branches), the MAP bit probability can be broken into the past (beginning of trellis to the present state), the present state (branch metric for the current value), and the future (end of trellis to current value). More specifically, the MAP decoder performs forward and backward recursions up to a present state, wherein the past and future probabilities are used along with the present branch metric to generate an output decision. The principles of providing soft output decisions are well known in the art, and several variations of the above-described decoding methods exist.

Most of the soft-input soft-output (SISO) decoders considered for turbo codes are based on the MAP algorithm in a paper by L.R. Bahl, J. Cocke, F. Jelinek, and J. Raviv entitled "Optimal Decoding of Linear Codes for Minimizing Symbol Error Rate", IEEE Transactions on Information Theory, Vol. IT-20, March 1974, pp. 284-7 (the "BCJR algorithm" or "BCJR method"). MAP algorithms not only minimize the probability of error for an information bit given the received sequence, they also provide the probability that the information bit is either a 1 or 0 given the received sequence. The BCJR decoding algorithm provides a soft output decision for each bit position (trellis section) wherein the influences of the soft inputs within the block are broken into contributions from the past (earlier soft inputs), the present soft

input, and the future (later soft inputs). The BCJR decoder algorithm requires a forward and a backward generalized Viterbi recursion on the trellis to arrive at an optimal soft output for each trellis section, or stage. These *a posteriori* probabilities, or more commonly the log-likelihood ratio (LLR) of the probabilities, are passed between SISO decoding steps in iterative turbo decoding. The LLR for information bit  $u_t$  is:

$$\Lambda_t = \log \frac{\sum_{(m,n) \in B^1} \alpha_{t-1}(n) \gamma_t(n,m) \beta_t(m)}{\sum_{(m,n) \in B^0} \alpha_{t-1}(n) \gamma_t(n,m) \beta_t(m)},$$

for all bits in the decoded sequence ( $t = 1$  to  $N$ ). In equation (1), the probability that the decoded bit is equal to 1 (or 0) in the trellis given the received sequence is composed of a product of terms due to the Markov property of the code. The Markov property may be stated as the assumption that the past and the future are independent given the present. The present,  $\gamma_t(n,m)$ , may be regarded as the probability of being in state  $m$  at time  $t$  and generating the symbol  $y_t$  when the previous state at time  $t-1$  was  $n$ . The present operates as a branch metric. The past,  $\alpha_t(m)$ , may be regarded as the probability of being in state  $m$  at time  $t$  with the received sequence  $\{y_1, \dots, y_t\}$ , and the future,  $\beta_t(m)$ , may be regarded as probability of generating the received sequence  $\{y_{t+1}, \dots, y_N\}$  from state  $m$  at time  $t$ . The probability  $\alpha_t(m)$  can be expressed as function of  $\alpha_{t-1}(m)$  and  $\gamma_t(n,m)$  and is called the forward recursion.

$$\alpha_t(m) = \sum_{n=0}^{M-1} \alpha_{t-1}(n) \gamma_t(n,m), \quad m = 0, \dots, M-1,$$

where  $M$  is the number of states. The reverse or backward recursion for computing the probability  $\beta_t(n)$  from  $\beta_{t+1}(n)$  and  $\gamma_t(n,m)$  is:

$$\beta_t(n) = \sum_{m=0}^{M-1} \beta_{t+1}(m) \gamma_t(n,m), \quad n = 0, \dots, M-1.$$

The overall *a posteriori* probabilities in equation (1) are computed by summing over the branches in the trellis  $B^1$  ( $B^0$ ) that correspond to  $u_t = 1$  (or 0).

The LLR in equation (1) requires both the forward and backward recursions to be available at time  $t$ . The BCJR method for meeting this requirement is to compute and store the entire backward recursion, and recursively compute  $\alpha_t(m)$  and  $\Lambda_t$  from  $t = 1$  to  $t = N$  using  $\alpha_{t-1}$  and  $\beta_t$ .

In terms of computational complexity, the BCJR method requires  $N*M$  state updates for the backward recursion ( $M$  state updates per trellis section,  $N$  trellis sections in the code) and provides optimal performance. In practice, a backward recursion is first performed by a processor across the entire block and stored in memory. The processor then performs a forward recursion. The results of the backward and forward recursions are used with the present state and stored future state to arrive at a soft output decision for each stage. In this case the processor operates on each state twice, once to generate and store the backward recursion states, and once to generate the forward recursion state.

To address the computational complexity and memory utilization problems of the soft output decoder, a sliding window method was developed. The sliding window technique is described in a paper by S. Benedetto, D. Divsalar, G. Montorsi, and F. Pollara, entitled "Algorithm for continuous

decoding of turbo codes,” Electronics Letters, Vol. 32, Feb. 15, 1996, pp. 314-315.

Another prior art decoder, described in U.S. Patent 5,933,462 to Viterbi et al. (and similarly in a paper of S. Pietrobon and S. Barbulescu, “A

5 Simplification of the Modified Bahl et al. Decoding Algorithm for Systematic Convolutional Codes,” Int. Symp. On Inform. Theory and its Applications, Sydney, Australia, pp. 1073-7, Nov. 1994, revised Jan. 4, 1996 and S. Pietrobon, “Efficient Implementation of Continuous MAP Decoders and a Synchronisation Technique for Turbo Decoders,” Int. Symp. On Inform.  
10 Theory and its Applications, Victoria, B.C., Canada, pp. 586-9, September 1996) comprises another sliding window technique.

The use of even the sub-optimal Max-Log-MAP algorithm for constituent code decoding makes heavy demands on processing resources. One prior art implementation of this algorithm, the Max-Log-MAP algorithm on the  
15 Motorola DSP56603 80 MIPS DSP, enables performance of 48.6 kbit/s. Given such a performance level, forty processors must work in parallel in order to provide the target real-time performance of 2Mbit/s, as defined in the 3G standards.

Another prior art device is the state-of-the-art Motorola StarCore SC140  
20 DSP, which cannot support a processing rate of more than 1 Mbit/s. A hand written assembly code required 36 cycles per code/iteration/bit, resulting in 288 cycles per 4 iterations, or equivalently ~1Mbit/s on a 300M cycles per second DSP.

## Summary of the Invention

It is a purpose of the present embodiments to provide a switchable data coding system that can be deployed in cellular wireless networks, to trade data rate with SINR when the conditions allow an increase in the SINR.

5 It is a purpose of the present embodiments to provide a data coding system that can be easily implemented in existing turbo encoders and decoders.

According to a first aspect of the present invention there is thus provided a switchable-output encoder for encoding an input data sequence to form an error protection encoded output sequence. The encoder is switchable between two encoding modes. The modes comprise a relatively complex mode suitable for a relatively high noise level channel and a relatively simple mode suitable for a relatively low noise level channel, wherein the relatively complex mode comprises a turbo coding mode.

10 In a preferred embodiment, the relatively simple mode comprises a degenerated version of the relatively complex mode.

In a further preferred embodiment, the relatively simple mode comprises a degenerated turbo coding mode.

In another preferred embodiment, the relatively simple mode comprises a convolutional coding mode.

20 In a further preferred embodiment, in the turbo coding mode, the output sequence comprises a multiplexed sequence containing at least three sub-sequences. The sub-sequences including a data sequence, a first coded sequence formable by encoding the data sequence, and a second coded

sequence formable by interleaving the data sequence into an interleaved sequence and encoding the interleaved sequence.

In a further preferred embodiment, in the degenerated turbo coding mode, the output sequence comprises a multiplexed sequence containing at least three sub-sequences. The sub-sequences include a data sequence, a first coded sequence formable by encoding the data sequence, and an interleaved sequence formable by interleaving the data sequence.

In another preferred embodiment, the switchable-output encoder comprises a first sub-encoder, to encode the input data sequence into a first coded sequence.

In a further preferred embodiment, the encoder comprises an interleaver, to interleave the input data sequence into an interleaved sequence.

In another preferred embodiment, the encoder further comprises a second sub-encoder, connected to the interleaver, to encode the interleaved data sequence into a second coded sequence.

In a preferred embodiment, the encoder further comprises a switch connected to the interleaver and to the second sub-encoder, wherein the switch is operable to provide one of the interleaved sequence and the second coded sequence as a switch output sequence, thereby affecting the composition of the encoder output sequence.

In a further preferred embodiment, the encoder further comprises an automatic controller, connected to the switch, the automatic controller being operable to monitor predetermined communication parameters in order to



determine a required one of the encoder modes, and to control switch operation accordingly.

In a preferred embodiment, in order to provide the turbo coding mode, the switch is settable to send the second coded sequence for output. In order to provide the degenerated turbo coding mode, the switch is settable to send the interleaved sequence for output.

In a further preferred embodiment, the encoder further comprises a multiplexer, connected to the encoder input, to the first sub-encoder, and to the switch, to multiplex the input data sequence, the first encoded sequence, and the switch output sequence into a single multiplexed sequence.

In another preferred embodiment, the multiplexed sequence serves as the error-protection encoded output sequence.

In a preferred embodiment, the first sub-encoder comprises a convolutional coder.

In a further preferred embodiment, the second sub-encoder comprises a convolutional coder.

In another preferred embodiment, the first sub-encoder and the second sub-encoder are recursive systematic convolutional encoders.

According to a second aspect of the present invention there is thus provided a switchable decoder for decoding a received sequence comprising error-protection encoded data, received from a noisy channel into an estimate of an input sequence, wherein the decoder is switchable between two modes, the modes comprising a relatively complex decoding mode suitable for a

relatively high noise level channel and a relatively simple decoding mode suitable for a relatively low noise level channel, and wherein the relatively complex mode comprises a turbo decoding mode.

5 In a preferred embodiment, the relatively simple decoding mode comprises a degenerated version of the relatively complex decoding mode.

In a further preferred embodiment, the relatively simple decoding mode comprises a degenerated turbo decoding mode.

In another preferred embodiment, the relatively simple decoding mode comprises a convolutional decoding mode.

10 In a preferred embodiment the decoder is operable to process the received sequence as a multiplexed sequence comprising at least three component sub-sequences.

15 In a preferred embodiment, when the decoder is in degenerated turbo decoding mode, the decoder is operable to process the first sub-sequence as a data sequence, the second sub-sequence as a directly encoded sub-sequence, and the third sub-sequence as an interleaved data sub-sequence.

In a preferred embodiment, the decoder comprises a separator, operable to separate the received data sequence into a first, a second, and a third data sub-sequence.

20 In a preferred embodiment, the decoder further comprises a first switch, connected to the sub-decoders, wherein the first switch is operable to connect the decoder output to the first sub-decoder output when the decoder is in relatively complex decoding mode, and to connect the decoder output to the

second sub-decoder output when the decoder is in relatively simple decoding mode.

In a further preferred embodiment, the first sub-decoder is operable as a turbo decoder, and the second sub-decoder is operable as a degenerated turbo decoder.

In another preferred embodiment, the degenerated turbo decoder comprises a de-interleaver for de-interleaving the third sub-sequence to form a de-interleaved sub-sequence.

In a further preferred embodiment, the degenerated turbo decoder further comprises a convolutional code decoder for decoding the first sub-sequence, the second sub-sequence, and the de-interleaved sub-sequence into the estimate of an input sequence.

In a preferred embodiment, the convolutional code decoder comprises a hard-decision trellis decoder.

In another preferred embodiment, the convolutional code decoder comprises a soft-decision trellis decoder.

In a preferred embodiment, the decoder further comprises a second switch, connected to the separator, wherein when the decoder is in relatively complex decoding mode the second switch is settable to connect the separator output sub-sequences to inputs of the first sub-decoder, and when the decoder is in relatively simple decoding mode the second switch is settable to connect the separator outputs to inputs of the second sub-decoder.

In a preferred embodiment, the decoder further comprises an automatic controller connected to the first switch, the automatic controller being operable to monitor predetermined communication parameters in order to determine a required one of the decoder modes, and to control switch operation

5 accordingly.

In a preferred embodiment, the decoder further comprises an automatic controller, connected to the second switch, the automatic controller being operable to monitor predetermined communication parameters in order to determine a required one of the decoder modes, and to control switch operation

10 accordingly.

According to a third aspect of the present invention there is thus provided a switchable data encoder-decoder system, comprising a switchable-output encoder for encoding an input sequence to form an error protection encoded output sequence and a switchable decoder, for decoding a received

15 sequence into an estimate of the input sequence, wherein the encoder and the decoder are synchronously switchable between two modes of operation, the modes comprising a relatively complex mode suitable for a relatively high noise level channel and a relatively simple mode suitable for a relatively low noise level channel, and wherein the relatively complex mode comprises a

20 turbo coding/decoding mode.

In a preferred embodiment, the relatively simple mode comprises a degenerated version of the relatively complex mode.

In a further preferred embodiment, the relatively simple mode comprises a degenerated turbo coding/decoding mode.

In a further preferred embodiment, the relatively simple mode comprises a convolutional coding/decoding mode.

5 In another preferred embodiment, when the encoder-decoder system is in turbo coding/decoding mode the encoder is operable to output a multiplexed signal comprising three sub-sequences, the sub-sequences comprising the input data sequence, a first coded sequence, and an interleaved and encoded data sequence.

10 In a further preferred embodiment, when the encoder-decoder system is in degenerated turbo coding/decoding mode the encoder is operable to output a multiplexed signal comprising three sub-sequences, the sub-sequences comprising the input data sequence, a first coded sequence, and an interleaved data sequence.

15 In a further preferred embodiment, when the encoder-decoder system is in degenerated turbo coding/decoding mode the decoder is operable to decode a received version of a multiplexed signal comprising the input data sequence, a first coded sequence, and an interleaved data sequence into an estimate of the input sequence.

20 In a preferred embodiment, the encoder comprises: an interleaver, a first sub-encoder, a second sub-encoder connected to the interleaver, a switch connected to the interleaver and to the second sub-encoder, and a multiplexer connected to the encoder input, the first sub-encoder, and the switch. The

interleaver interleaves the input signal into an interleaved data sequence. The first sub-encoder encodes the input sequence into a first coded sequence. The second sub-encoder encodes the input sequence into a second coded sequence. The switch is settable to provide the second coded sequence as a switch output sequence when the system is in turbo coding/decoding mode, and to provide the interleaved data sequence as a switch output sequence when the system is in degenerated turbo coding/decoding mode. And the multiplexer multiplexes the data sequence, the first coded sequence, and the switch output sequence into an output sequence

In a further preferred embodiment, the decoder comprises: a separator, a first sub-decoder connected to the separator, a second sub-decoder connected to the separator, a first switch connected to the sub-decoders, and a second switch connected between the separator and the sub-decoders. The separator separates the received data sequence into a first, a second, and a third data sub-sequence.

The first sub-decoder decodes the sub-sequences when the encoder-decoder system is in relatively complex mode. The second sub-decoder decodes the sub-sequences when the encoder-decoder system is in relatively simple mode.

The first switch connects the decoder output to the first sub-decoder output when the decoder is in relatively complex decoding mode, and connects the

decoder output to the second sub-decoder output when the decoder is in relatively simple decoding mode. And the second switch routes the sub-sequences to either of the first and second sub-decoders in accordance with a current mode of operation.

In a further preferred embodiment, the encoder-decoder system further comprises an automatic controller, connected to the encoder switch and to the decoder first switch, the automatic controller being operable to monitor predetermined communication parameters in order to determine a required one of the encoder-decoder system modes, and to control switch operation accordingly.

In a further preferred embodiment, the first sub-decoder comprises a turbo code decoder.

In another preferred embodiment, the second sub-decoder comprises: a de-interleaver, connected to the separator, for de-interleaving the third sub-sequence to form a de-interleaved sub-sequence, and a convolutional code decoder, connected to the separator and to the de-interleaver, for decoding the first sub-sequence, the second sub-sequence, and the de-interleaved sub-sequence into the estimate of an input sequence.

In a further preferred embodiment, the encoder-decoder system further comprises an automatic controller, connected to the second switch, the automatic controller being operable to monitor predetermined communication parameters in order to determine a required one of the decoder modes, and to control switch operation accordingly.

According to a fourth aspect of the present invention there is thus provided a method for encoding an input data sequence into an error protection encoded output sequence, comprising: receiving an input data sequence, interleaving the input sequence to form an interleaved data sequence, encoding

the input sequence to form a first encoded sequence according to a first coding rule, encoding the interleaved sequence to form a second encoded sequence according to a second coding rule, selecting either one of the interleaved and the second encoded sequence, and multiplexing the input sequence, the first  
5 encoded sequence, and the selected sequence to form the error protection encoded output sequence.

In a preferred embodiment, the selection is made based on current values of predetermined communication parameters.

In a further preferred embodiment, the first encoding rule comprises  
10 convolutional coding.

In a further preferred embodiment, the second encoding rule comprises convolutional coding.

According to a fifth aspect of the present invention there is thus provided a method for decoding a received sequence comprising error-  
15 protection encoded data received from a noisy channel into an estimate of an input sequence, comprising: receiving the sequence from the noisy channel, separating the received sequence into a first, a second, and a third data sub-sequence, selecting either one of a first sub-decoder and a second sub-decoder, and decoding the sub-sequences into the estimate of an input sequence using  
20 the selected sub-decoder.

In a preferred embodiment, selection is made based on current values of predetermined communication parameters.



In a further preferred embodiment, the first sub-decoder comprises a turbo code decoder.

In a further preferred embodiment, the method by which the second sub-decoder decodes the first, second, and third data sub-sequences comprises: de-  
5 interleaving the third sub-sequence into a deinterleaved sub-sequence, and decoding the first, the second, and the de-interleaved sub-sequences into the estimate of an input sequence using a convolutional code decoder.

In a further preferred embodiment, the convolutional code decoder comprises a hard-decision trellis decoder.

10 In another preferred embodiment, the convolutional code decoder comprises a soft-decision trellis decoder.

### **Brief Description of the Drawings**

For a better understanding of the invention and to show how the same may be carried into effect, reference will now be made, purely by way of  
15 example, to the accompanying drawings, in which:

Figure 1 shows a simplified block diagram of a conventional turbo code encoder.

Figure 2 shows a simplified block diagram of a switchable output encoder according to a preferred embodiment of the present invention.

20 Figure 3 shows a simplified block diagram of a switchable output encoder according to a preferred embodiment of the present invention.

Figure 4 shows a simplified block diagram of a switchable decoder according to a preferred embodiment of the present invention.

Figure 5 shows a simplified block diagram of a simple sub-decoder according to a preferred embodiment of the present invention.

Figure 6 is a simplified flow chart of a method for encoding an input sequence according to a preferred embodiment of the present invention.

5 Figure 7 is a simplified flow chart of a method for decoding a received sequence containing error-protection encoded data according to a preferred embodiment of the present invention.

10 Figure 8 is a simplified flow chart of a method for use by a sub-decoder for decoding received sub-sequences according to a preferred embodiment of the present invention.

### **Description of the Preferred Embodiments**

15 Currently known iterative decoding algorithms are significantly less complex than optimal decoding, yet remain computationally complex due to the use of costly component decoders. While the number of system users and the channel characteristics vary over the lifetime of a system, existing systems provide only a single coding solution, which may be sub-optimal for the existing conditions. There is a need for a more flexible data coding system. The dual-mode coding system embodiments described below provide a solution that enables trading data rate with signal to interference noise ratio (SINR) 20 when conditions allow.

Reference is now made to Figure 1, which is a simplified block diagram of a conventional turbo code encoder 10. The encoder 10 comprises an interleaver 12, two recursive systematic convolutional (RSC) encoders 14 and

16, and a multiplexer 18. The input to the turbo code encoder 10 is an input data sequence  $\mathbf{d}$ . The first encoder  $\text{RSC}_0$  14 encodes the input data sequence  $\mathbf{d}$  directly, forming data sequence  $\mathbf{Y}^0$ . Interleaver 12 interleaves the input data sequence  $\mathbf{d}$ . The interleaved sequence is the input to the second encoder  $\text{RSC}_1$

5 16.  $\text{RSC}_1$  encoder 16 encodes the interleaved version of data sequence  $\mathbf{d}$ , forming data sequence  $\mathbf{Y}^1$ . The multiplexer 18 combines the three data sequences into output sequence  $\mathbf{C}$ . The coded sequence is given by  $\{\mathbf{C}_k\}$ , with  $\mathbf{C}_k = (\mathbf{X}_k, \mathbf{Y}_k^0, \mathbf{Y}_k^1)$ , where  $\mathbf{X}_k = \mathbf{d}_k$  is the information bit at time  $k$ ,  $\mathbf{Y}_k^0$  is the parity output bit from the encoder  $\text{RSC}_0$  14, and  $\mathbf{Y}_k^1$  is the parity output bit  
10 from the encoder  $\text{RSC}_1$  16. This output sequence is then transmitted over a noisy channel. For notational simplicity, when reference is made to an entire sequence, the  $k$  subscript, used to indicate a specific bit in the sequence, will be omitted.

Reference is now made to Figure 2, which is a simplified block diagram  
15 of a switchable output encoder 30 according to a first embodiment of the present invention. The encoder 30 encodes an input data sequence to form an error-protection encoded output sequence, for transmission over a noisy channel. Encoder 30 operates in two encoding modes: a relatively complex mode, suitable for a relatively high noise level channel, and a relatively simple  
20 mode, suitable for a relatively low noise level channel. The relatively complex mode comprises a turbo coding mode. In a preferred embodiment, the relatively simple mode is a degenerated form of the relatively complex mode. The encoder 30 comprises interleaver 32, sub-encoder A 34, sub-encoder B 36,

and multiplexer 38, all of which operate in a manner similar to those of the encoder of Figure 1. The encoder 30 additionally comprises switch A 40, which determines the composition of the encoder output according to a selected encoding mode.

5 The input to the switchable-output encoder 30 is an input data sequence  $\mathbf{d}^0$ . Sub-encoder A 34 encodes the input data sequence  $\mathbf{d}^0$  directly, forming data sequence  $\mathbf{d}^1$ . Interleaver 12 interleaves the input data sequence  $\mathbf{d}^0$ , to form  $\mathbf{d}^{\text{int}}$ . This interleaved sequence  $\mathbf{d}^{\text{int}}$  is the input to sub-encoder B 36. Sub-encoder B 36 encodes  $\mathbf{d}^{\text{int}}$ , forming data sequence  $\mathbf{d}^2$ . Unlike the turbo encoder described above, data sequence  $\mathbf{d}^2$  is not input directly to one of the multiplexer inputs. Sequence  $\mathbf{d}^2$  serves as one of the inputs to switch A 40. The second input to switch A 40 is the interleaved sequence  $\mathbf{d}^{\text{int}}$ . The switch output sequence  $\mathbf{d}^{\text{sw}}$  is selected according to the current encoding mode. When the encoder 30 is in relatively complex mode, switch A 40 is set to position **a**, and the switch output sequence  $\mathbf{d}^{\text{sw}}$  is sequence  $\mathbf{d}^2$ . When the encoder 30 is in relatively simple mode, sub-encoder 36 B is bypassed. In this mode, switch A 40 is set to position **b**, and the switch output sequence  $\mathbf{d}^{\text{sw}}$  is sequence  $\mathbf{d}^{\text{int}}$ . The multiplexer 38 combines the three data sequences,  $\mathbf{d}^0$ ,  $\mathbf{d}^1$ , and  $\mathbf{d}^{\text{sw}}$ , into output sequence  $\mathbf{d}^{\text{OUT}}$ . The coded sequence is given by  $\{d_k^{\text{OUT}}\}$ , with

20  $d_k^{\text{OUT}} = (d_k^0, d_k^1, d_k^{\text{sw}})$ , where  $d_k^0$  is the information bit at time  $k$ ,  $d_k^1$  is the parity output bit from sub-encoder A 34, and  $d_k^{\text{sw}}$  is the parity output bit from switch A 40.

In the preferred embodiment of the encoder, sub-encoder A 34 and sub-encoder B 36 are recursive systematic convolutional (RSC) encoders. The relatively complex encoder mode is equivalent to conventional turbo encoding, and the relatively simple mode is a degenerated turbo coding mode.

5 In an alternate preferred embodiment, the relatively simple mode comprises a convolutional coding mode.

Reference is now made to Figure 3, which is a simplified block diagram of a further preferred embodiment of a switchable output encoder 50. This embodiment is similar to the embodiment of Figure 2, with the addition of switch B 62 introduced after interleaver 52. Switch B 62 is synchronized with switch A 60, so that the interleaver output sequence  $\mathbf{d}^{\text{int}}$  is connected to encoder B 56 only when the encoder 50 is in relatively complex mode. When the encoder 50 is in relatively simple mode,  $\mathbf{d}^{\text{int}}$  is connected to switch A 60, thereby eliminating an unnecessary encoding of the interleaved sequence.

15 Reference is now made to Figure 4, which is a simplified block diagram of a preferred embodiment of a switchable decoder 90. Switchable decoder 90 decodes a received sequence  $\mathbf{R}$  containing error-protection encoded data received from a noisy channel, into an estimate of an input sequence. The decoder operates in two decoding modes corresponding to the two  
20 aforementioned encoding modes: a relatively complex mode suitable for a relatively high noise level channel and a relatively simple mode suitable for a relatively low noise level channel. In a preferred embodiment, the relatively simple mode is a degenerated version of the relatively complex mode.

Switchable decoder 90 comprises separator 92, complex sub-decoder 94, simple sub-decoder 96, a switch A 98, and a switch B 100. Switch B 100 is synchronized with switch A 98, so that the separator output sub-sequences are connected to complex sub-decoder 94 only when the decoder 90 is in relatively complex mode. When the decoder 90 is in relatively simple mode, the separator output sub-sequences are connected to simple sub-decoder 96.

Separator 92 processes a multiplexed received data sequence, to separate the multiplexed sequence into its component sub-sequences,  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^2$ . The component sub-sequences are then decoded by complex sub-decoder 94, to form decoded sequence  $\hat{\mathbf{d}}^{comp}$ , or by simple sub-decoder 96, to form decoded sequence  $\hat{\mathbf{d}}^{simp}$ . The decoded sequence is the input to switch A 98, which selects an output sequence  $\hat{\mathbf{d}}^0$  according to the current decoding mode. When the decoder 90 is in relatively complex mode, switches A 98 and B 100 are set to position **a**, and the switch output sequence  $\hat{\mathbf{d}}^0$  is thus formed by sequence  $\hat{\mathbf{d}}^{comp}$ . When the decoder 90 is in relatively simple mode, switches A 98 and B 100 are set to position **b**, and the switch output sequence  $\hat{\mathbf{d}}^0$  is formed by sequence  $\hat{\mathbf{d}}^{simp}$ .

Specific embodiments of the sub-decoders are determined by the coding method used prior to transmission. In a preferred embodiment of a complex sub-decoder, the received data sequence input to the decoder is turbo code encoded, and the complex sub-decoder is a turbo code decoder known in the art.

Reference is now made to Figure 5, which is a simplified block diagram of a preferred embodiment of a simple sub-decoder 110. The inputs to the sub-decoder 110 are the sub-sequences  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^2$ . In the preferred embodiment, sub-decoder 110 is a component of the switchable decoder described above for Figure 4, and the sub-sequences are provided by a separator, as previously described. In the preferred embodiment, sub-sequence  $\mathbf{r}^0$  is a noisy version of the input signal. Sub-sequence  $\mathbf{r}^1$  is a noisy version of a sub-sequence formed by encoding the input signal. In a preferred embodiment the encoding method used to form sub-sequence  $\mathbf{r}^1$  is convolutional coding. Sub-sequence  $\mathbf{r}^2$  is a noisy version of a sub-sequence formed by interleaving the input signal. First, sub-sequence  $\mathbf{r}^2$  is de-interleaved by de-interleaver 112, to form  $\mathbf{r}^{\text{dint}}$ . Like sequence  $\mathbf{r}^0$ , sequence  $\mathbf{r}^{\text{dint}}$  is a noisy version of the input signal. After de-interleaving, the values  $(\mathbf{r}_k^0, \mathbf{r}_k^1, \mathbf{r}_k^{\text{dint}})$  are all associated with the k-th bit of the input sequence. Since  $\mathbf{r}_k^0$  and  $\mathbf{r}_k^{\text{dint}}$  were received at different times, their noise content is statistically independent. The de-interleaved sequence  $\mathbf{r}^{\text{dint}}$  thereby introduces new information about the input sequence, which can be used by sub-decoder 110 during the decoding process. The sequences  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^{\text{dint}}$  are input into a convolutional code decoder 114, which decodes them into an estimate of the input sequence  $\hat{\mathbf{d}}^{\text{simp}}$ . In one preferred embodiment, convolutional code decoder 114 is a non-iterative soft-decision trellis decoder known in the art. In another preferred embodiment, convolutional code decoder 114 is a non-iterative hard-decision trellis decoder known in the art.

Sub-sequences  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^{\text{dint}}$  are used to calculate the convolutional code decoder metrics.

In a preferred embodiment, the simple sub-decoder is a component of the complex sub-decoder.

5 Preferably, the switchable-output encoder and the switchable decoder operate as a system. In such a system, the switches in the encoder and decoder are synchronized, so that both the encoder and the decoder operate in the same mode. The complex sub-decoder decodes output sequences produced by the encoder during complex mode operation, and the simple sub-decoder decodes  
10 output sequences produced by the encoder during simple mode operation.

Reference is now made to Figure 6, which is a simplified flow chart of a method for encoding an input sequence according to a preferred embodiment of the present invention. The input data sequence  $\mathbf{d}^0$  is received in step 132. An interleaved sequence,  $\mathbf{d}^{\text{int}}$ , is formed in step 134 by interleaving  $\mathbf{d}^0$ . In step 136,  
15 sequence  $\mathbf{d}^0$  is encoded to form sequence  $\mathbf{d}^1$ . In step 138, the interleaved sequence  $\mathbf{d}^{\text{int}}$  is encoded to form an additional sequence  $\mathbf{d}^2$ . The encoding rule used to form sequence  $\mathbf{d}^2$  may differ from the encoding rule used to form sequence  $\mathbf{d}^1$ . One of the sequences  $\mathbf{d}^{\text{int}}$  or  $\mathbf{d}^2$  is selected in step 140. Finally, in step 142, an output sequence  $\mathbf{d}^{\text{OUT}}$  is formed by multiplexing  $\mathbf{d}^0$ ,  $\mathbf{d}^1$ , and the  
20 selected sequence into a single error-protection encoded sequence.

Reference is now made to Figure 7, which is a simplified flow chart of a method for decoding a received sequence containing error-protection encoded data according to a preferred embodiment of the present invention. Received



signal  $\mathbf{R}$  is input to the system in step 152. In step 154,  $\mathbf{R}$  is separated into component sub-sequences  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^2$ . A sub-decoder is selected in step 156. The selected decoder decodes the sub-sequences to form an estimate of an input sequence,  $\hat{d}^0$ , in step 158. This estimated sequence serves as the decoder output.

Reference is now made to Figure 8, which is a simplified flow chart of an embodiment of a method suitable for use by a sub-decoder for decoding received sub-sequences. First, sub-sequences  $\mathbf{r}^0$ ,  $\mathbf{r}^1$ , and  $\mathbf{r}^2$  are received in step 172. In a preferred embodiment, these sequences correspond to those of the method of Figure 7 described above. These sub-sequences are decoded in step 174 by a convolutional code decoder into an estimate of an original data sequence  $\hat{d}^{simp}$ .

When deploying cellular wireless networks, systems are often designed to meet certain quality of service (QoS) requirements by employing sophisticated and computationally demanding error correction schemes. During the early stages of deployment there may be only a small number of users, and consequently lower levels of interference. Nevertheless, even with only a few customers in a system (i.e., lower interference levels), complex error correction is still provided at the handsets. Alternatively, standard DSPs may be used, resulting in lower peak data rates. The present embodiments can be employed to ease this problem, typically encountered in the early stages of network deployment, by allowing an increase in the data rate without sacrificing the QoS.

In later stages of network deployment, the present embodiments may be used for trading data rate with SINR when the conditions allow an increase in the SINR. For example, when the user is near the base station the power transmitted by the base station to that particular user is relatively low, and can be increased without significantly affecting other users in the cell or neighboring cells.

The present embodiments provide a switchable coding system which can operate over a variety of conditions. The relatively complex mode provides a robust solution for networks operating with a high number of users, or under poor SINR conditions. The relatively simple mode is suitable for systems operating with low numbers of users, or with high SINR. The relatively simple coding mode provides a simpler coding and decoding technique, requiring less processing and enabling a higher data rate. The current system mode of operation is preferably selected automatically, according to detected network and channel conditions.

It is appreciated that certain features of the invention, which are, for clarity, described in the context of separate embodiments, may also be provided in combination in a single embodiment. Conversely, various features of the invention which are, for brevity, described in the context of a single embodiment, may also be provided separately or in any suitable subcombination.

It will be appreciated by persons skilled in the art that the present invention is not limited to what has been particularly shown and described

hereinabove. Rather the scope of the present invention is defined by the  
appended claims and includes both combinations and subcombinations of the  
various features described hereinabove as well as variations and modifications  
thereof which would occur to persons skilled in the art upon reading the  
5 foregoing description.